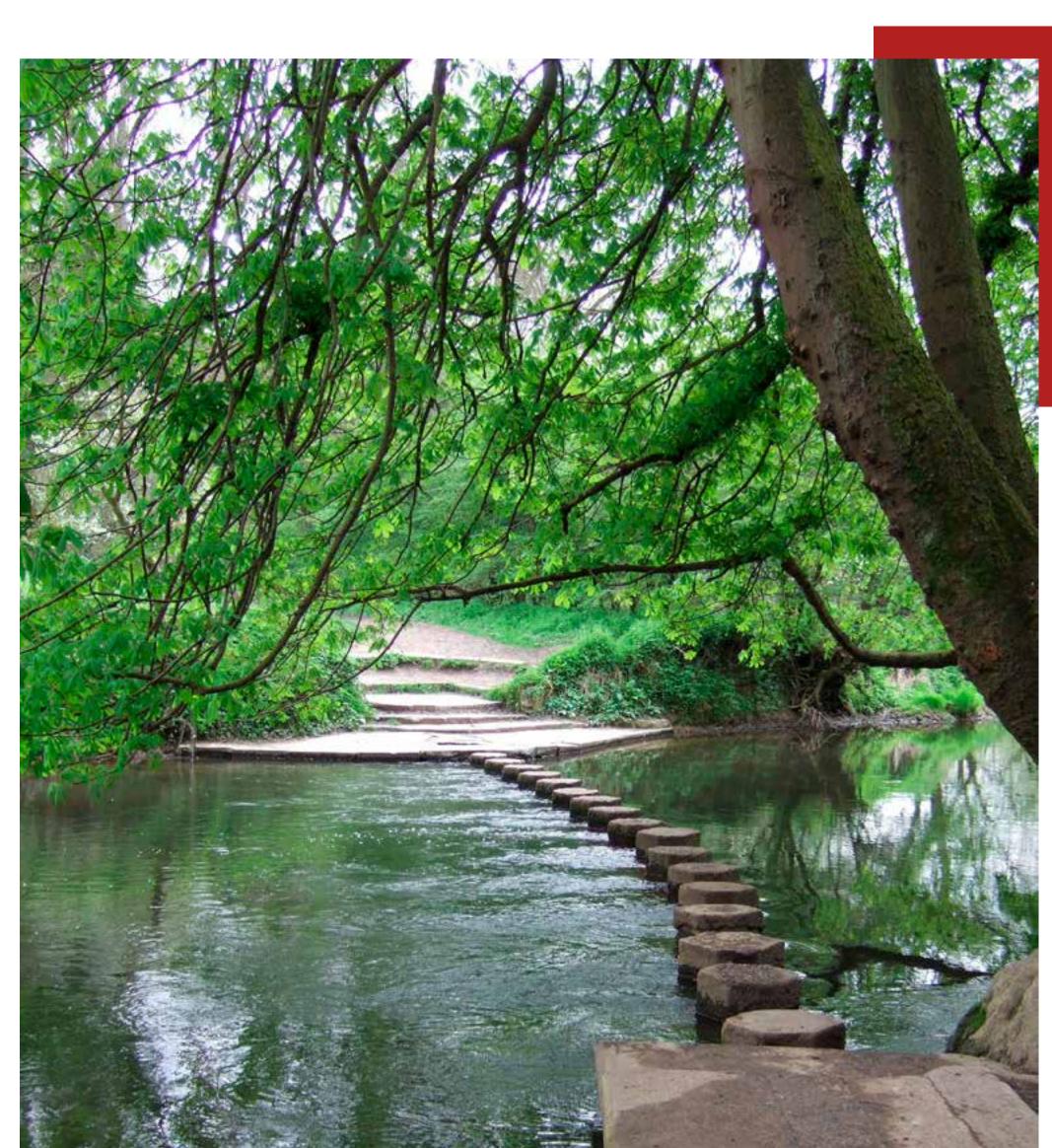
## $\vdash \exists x.\varphi_1(x), \exists y.\varphi_1'(x,y), \exists z.\varphi_2'(x,y,z)$ $-\exists x.\varphi_1(x), \varphi_1(x), \forall y\exists z.\varphi_2'(x,y,z),$ Machine

 $\vdash \Gamma, x < y, \neg x < y, x < z \land z < y$ 

# $-\exists x.\varphi_1(x), \varphi_1(x), \exists x.\varphi_2(x)$ Checked Proofs $\frac{\exists x.\varphi_1(x),\exists x.\varphi_2(x)}{\vdash (\exists x.\varphi_1(x))\lor(\exists x.\varphi_2(x))}$

## When computers improve mathematical rigour

Since the invention of the concept of proof in ancient Greece, mathematicians have always sought to write ever more rigourous proofs: identifying axioms precisely, defining every object used in the proof, avoiding the call to intuition, etc. Machine-checked proof is a new step in this never ending quest of rigour. A machine-checked proof is written with such precision that a computer program can check its correctness.



Like the crossing of a river ford, a mathematical proof goes step by step

#### Goal ((P -> Q) -> P) -> P. intro piqip. assert (ponp: $P \setminus / \sim P$ ). exact (classic P). destruct ponp as [p|np]. assumption. apply piqip. intro p. destruct np. assumption. Qed.

```
\Gamma, \neg P, P, \neg Q \vdash \neg P
                                                                                                                                                                             \Gamma, \neg P, P, \neg Q \vdash P
                                            \Gamma, \neg P \vdash (P \rightarrow Q) \rightarrow P
\overline{\Gamma, \neg P \vdash \neg P}^{\text{ (ax)}}
                                                                                                     \Gamma, \neg P \vdash P
```

 $-\Gamma, x < y \land (\neg x < z \lor \neg z < y), \neg x < y$ 

 $\vdash \Gamma, x < y \land (\neg x < z \lor \neg z < y), \neg x < y \lor$ 

 $\vdash \exists x. \varphi_1(x), \exists y. \varphi_1'(x, y), x < y \land (\neg x < z \lor \neg z < y)$ 

 $\vdash \exists x. \varphi_1(x), \exists y. \varphi_1'(x, y), \forall z. x < y \land (\neg x < z \lor \neg z < y)$ 

 $\vdash \Gamma$ , -

Two proofs of Peirce's law, in COQ (above) and in the natural deduction calculus (below).

### THE BEGINNING

The two first proof-checkers were Automath (de Bruijn, 1967), and then LCF (Milner, 1972). Their goals were different: Automath was designed to check general mathematical proofs, LCF, more specifically, proofs of properties of programs.

## TODAY

The development of proof-checkers triggered the development of new theories, besides set theory, to express mathematics: each system innovates, introducing new features to express mathematical statements and proofs, just like each new programming language introduces new features to express programs.

Popular proof-checkers are ACL 2, Agda, Coq, HOL Light, HOL 4, Lean, Mizar, Nuprl, PVS, and many others. These proof-checkers are specific to one theory. Others, such as Beluga, Dedukti, Isabelle, Lambdaprolog, Twelf, and others are frameworks, where various theories can be defined.

They have in total more than 10,000 users.

### RECENT PROOFS

2000 : four colour theorem (Gonthier et al.)

2008: correctness of the C compiler CompCert (Leroy et al.) 2009: correctness of the operating system seL4 (Klein et al.)

2012 : Feit-Thompson theorem (Gonthier et al.) 2014 : Kepler's conjecture (Hales et al.)

2014: UniMath a body of mathematics using univalent foundations (Voevodsky et al.)

Several of these projects aim at gathering a substantial body of mathematics, like Euclid's *Elements* and Bourbaki's *Élements de* mathématique did.



For long, mathematics was the only science not to use instruments. The computer is becoming the telescope of mathematicians